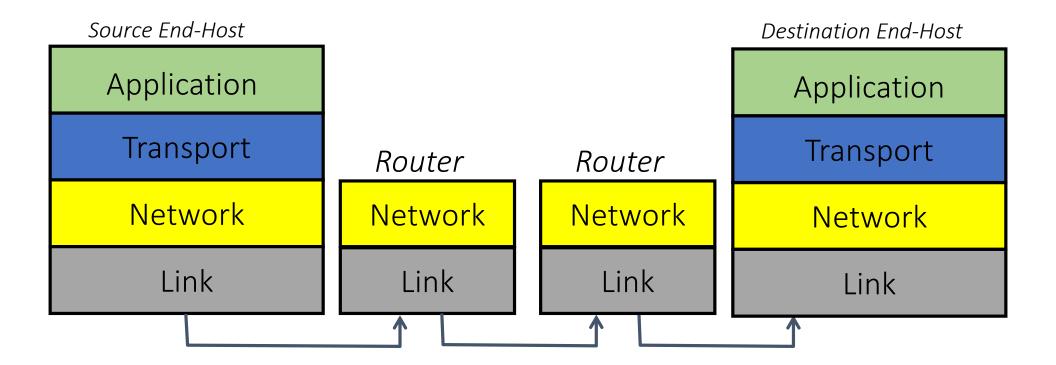
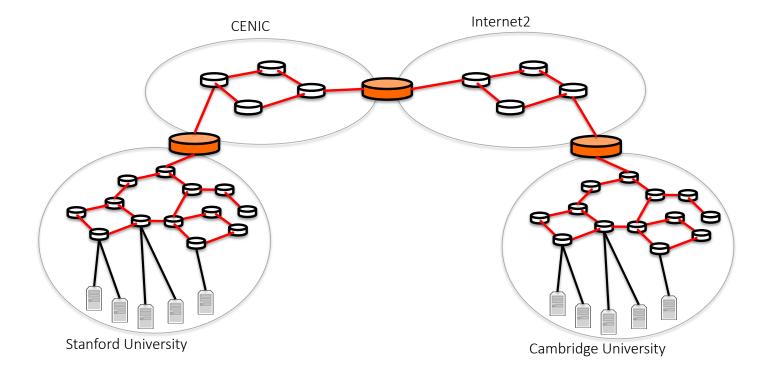
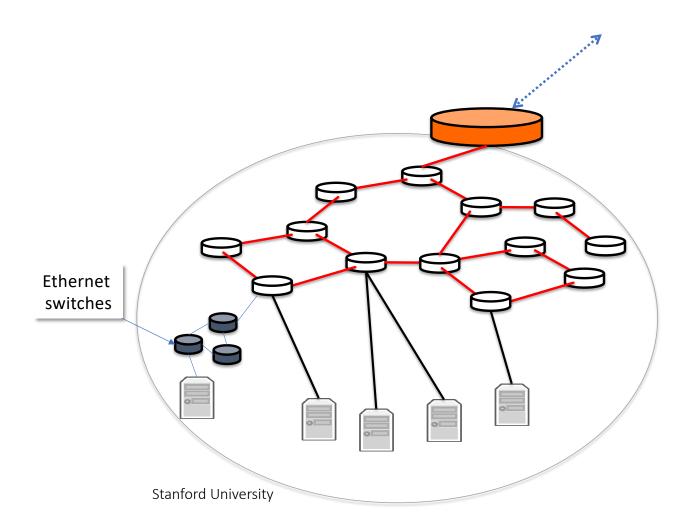
Ethernet and CSMA/CD

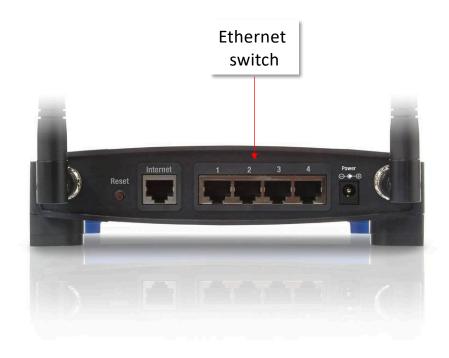
The 4 Layer Internet Model



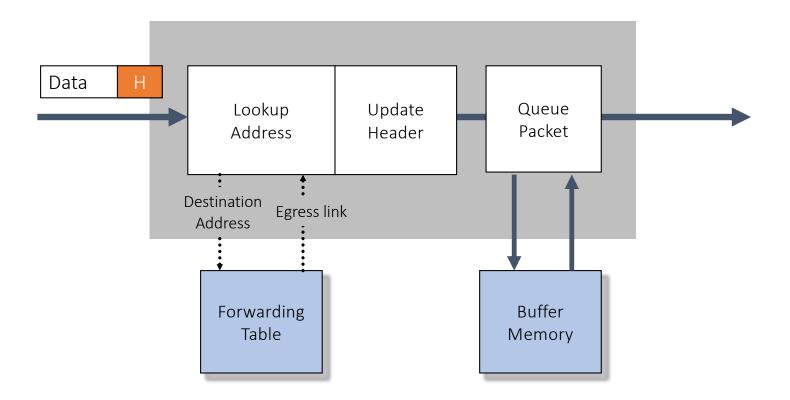




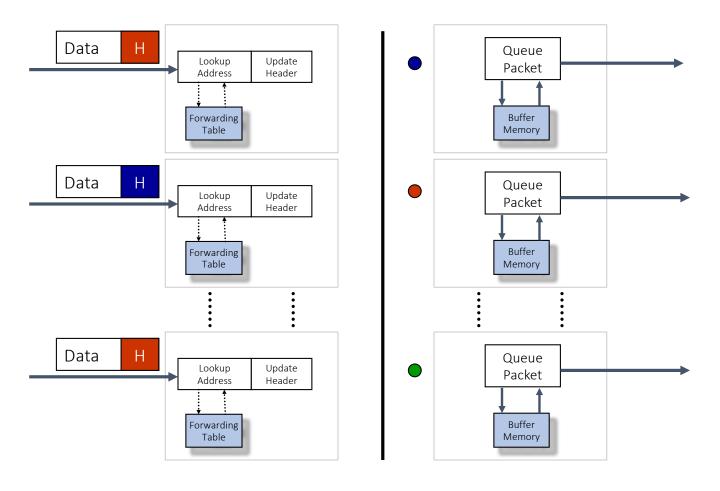




Generic Packet Switch



Generic Packet Switch



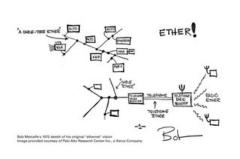
Ethernet Switch

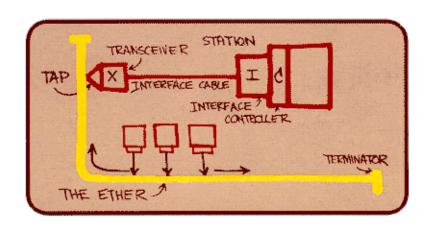
- 1. Examine the header of each arriving frame.
- 2. If the Ethernet DA (aka "MAC Address") is in the forwarding table, forward the frame to the correct output port(s).
- 3. If the Ethernet DA is not in the table, broadcast the frame to <u>all</u> ports (except the one through which the frame arrived).
- 4. Entries in the table are <u>learned</u> by examining the Ethernet SA of arriving packets.

Internet Router

- 1. If the Ethernet DA of the arriving frame belongs to the router, accept the frame. Else drop it.
- 2. Examine the IP version number and length of the datagram.
- 3. Decrement the TTL, update the IP header checksum.
- 4. Check to see if TTL == 0.
- 5. If the IP DA is in the forwarding table, forward to the correct egress port(s) for the next hop.
- 6. Find the Ethernet DA for the next hop router.
- 7. Create a new Ethernet frame and send it.

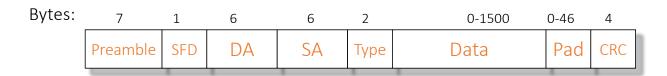
The Original Ethernet





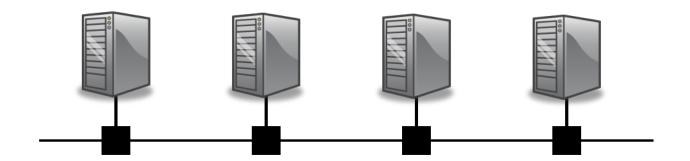
Original pictures drawn by Bob Metcalfe, co-inventor of Ethernet (1972 – Xerox PARC)

Ethernet Frame Format



- 1. Preamble: trains clock-recovery circuits
- 2. Start of Frame Delimiter: indicates start of frame
- **3. Destination Address**: 48-bit globally unique address assigned by manufacturer.
 - 1b: unicast/multicast
 - 1b: local/global address
- 4. Type: Indicates protocol of encapsulated data (e.g. IP = 0x0800)
- 5. Pad: Zeroes used to ensure minimum frame length
- 6. Cyclic Redundancy Check: check sequence to detect bit errors.

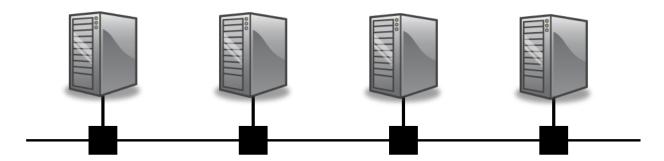
The origins of Ethernet



Sharing a "medium"

- Ethernet is (or at least was originally) an example of multiple hosts sharing a common cable ("medium").
- To share the medium, we need to decide who gets to send, and when.
- There is a general class of "Medium Access Control Protocols", or MAC Protocols.

CSMA/CD Protocol



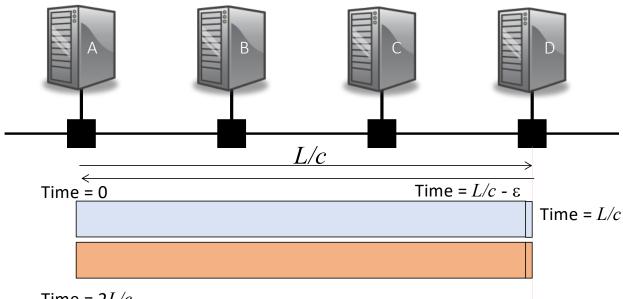
All hosts transmit & receive on one channel Packets are of variable size.

When a host has a packet to transmit:

- 1. Carrier Sense: Check the line is quiet before transmitting.
- 2. Collision Detection: Detect collision as soon as possible. If a collision is detected, stop transmitting; wait a <u>random time</u>, then return to step 1.

binary exponential backoff

CSMA/CD Packet size requirement

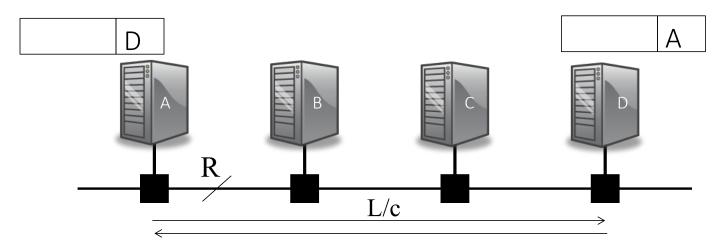


Time = 2L/cFirst time "A" knows about collision

Therefore, "A" is guaranteed to know about the collision while it is still transmitting if:

$$\frac{P}{R} > \frac{2I}{c}$$

CSMA/CD Min packet size requirement



For an end host to detect a collision before it finishes transmitting a packet, we require:

$$\frac{P}{R} > \frac{2L}{c}$$

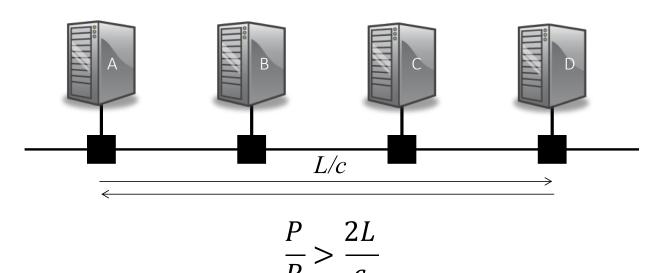
where P is the size of a packet.

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CSMA/CD Min packet size requirement

Example:

R = 10Mb/s, L = 10,000m, $c = 2 \times 10^8 m/s$.



$$\therefore P_{min} = \frac{2LR}{c} = \frac{2 \times 10^{11}}{2 \times 10^{8}} = 1,000 \ bits$$